

## Inter-process Communication

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## Potential Topics of Term Papers

The topic should be related to distributed systems in general, such as

- Group Communications,
- Peer-to-Peer systems,
- Overlay Networks,
- Grid Computing,
- Object Middleware,
- and others.

## Potential Topics of Term Papers

- ♦ Resource Discovery and Management
- ♦ Security and Policy Management
- ♦ Resource Scheduling and Load Balancing
- ♦ Synchronization (e.g., clock synchronization, Election Algorithm, mutual exclusion, etc.)
- ♦ Consistency and replication
- ♦ Reliability and Survivability
- ♦ Performance Evaluation
- ♦ Anonymity. Censor-resistant
- ♦ Workload characterization
- ♦ Multi-cast Fingerprinting
- ♦ Anonymous authentication in dynamic group communications
- ♦ Data Replication strategies for Grid Computing systems
- ♦ Reputation-based resource scheduling for Grid Computing systems
- ♦ Middleware-based application design and development (e.g., real-time CORBA, fault-tolerant CORBA, etc.)
- ♦ You are welcome to propose your own topic!!!

## Your Term Papers

- ♦ Every student is required to finish a term paper.
- ♦ Deadlines:
  - February 26 (Thursday, 5:00pm), topic selection due (UG).
  - March 11 (Thursday, 5:00pm), 3-pages proposal (problem definition) due (UG).
  - March 25 (Thursday, 5:00pm), 5-pages solutions and drafted evaluation plan due (G).
  - April 6 (Tuesday, 5:00pm), 5-pages summary of the papers you read and important issues you think (U)
  - April 15 (Thursday 5:00pm), experimental results and improved solution due (G).
  - April 27 (Tuesday 5:00pm), 10-page literature survey (U) and 15-pages full term paper (G) due, including problem definition, solutions, experimental data, conclusion and future works.
- ♦ Paper format: Latex or WORD, IEEE transactions, please refer: <http://www.ieee.org/organizations/pubs/transactions/stylesheets.htm>

## Readings for Today's Lecture

- References
  - Chapter 2 of "Distributed Systems: Principles and Paradigms"

## Object-Oriented Distributed Technology

- ♦ Objects
- ♦ Objects in Distributed Systems
- ♦ Requirements of Multi-User Applications

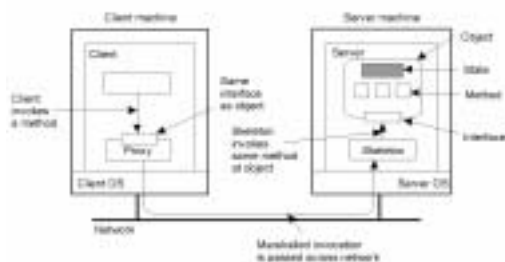
## Object-Oriented Languages

- ◆ Object Identity
  - "object identifiers" (OIDs)
  - OIDs as first class values
- ◆ Actions
  - Initiated by sending message to object requesting method invocation
  - State in object may change
  - cascaded invocations of methods
- ◆ Dynamic Binding
  - The method executed is chosen according to the class of the recipient of the message.
- ◆ Garbage Collection
  - Dynamically allocated instances may be explicitly *deleted* or space is freed implicitly by garbage collector.

## Objects in Distributed Systems

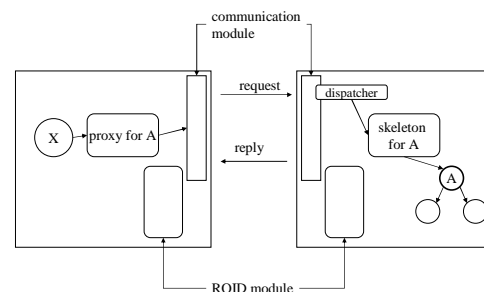
- ◆ Object Identity in a Distributed System
  - Remote object identifiers (ROIDs)
  - Ex. Java:  $ROID = \text{endpoint (Java vm)} + \text{identifier (ObjID)}$
  - ROIDs as first-class values
  - Service for comparing remote object identifiers
  - e.g. Java: *RemoteObject::equals()*
- ◆ Actions in a Distributed Object System
  - Remote Method Invocation
- ◆ The Role of Proxies for Transparent RMI
  - Local proxy for each remote object that can be invoked by local object.
  - Local proxy behaves like local object, but, instead of executing message, forwards it to the remote object. (client stubs)
  - Remote object has skeleton object with server stub procedures

## Distributed Objects



- ◆ Common organization of a remote object with client-side proxy.

## Proxies and Skeletons



## Proxies and Skeletons (cont)

- ◆ Proxies:
  - Need proxies to invoke remote objects.
  - Proxies are created when needed whenever ROID arrives in Reply message.
  - ROID module manages proxies and ROIDs.
- ◆ Dispatchers and Skeletons:
  - Not necessary for systems with reflection capabilities.
  - e.g. class *Method* in Java 1.2 reflection package: method *invoke* can be called on instance of *Method*.
  - Dispatcher now generic and skeleton unnecessary.

## Binding a Client to an Object

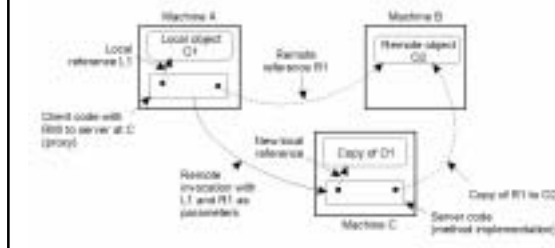
```

Distr_object* obj_ref;           //Declare a systemwide object reference
obj_ref = ...;                   // Initialize the reference to a distributed object
obj_ref-> do_something();         // Implicitly bind and invoke a method
                                (a)

Distr_object objPref;            //Declare a systemwide object reference
Local_object* obj_ptr;          //Declare a pointer to local objects
obj_ref = ...;                   //Initialize the reference to a distributed object
obj_ptr = bind(obj_ref);         //Explicitly bind and obtain a pointer to the local proxy
obj_ptr-> do_something();         //Invoke a method on the local proxy
                                (b)
  
```

- (a) Example with implicit binding using only global references
- (b) Example with explicit binding using global and local references

## Parameter Passing



- ◆ The situation when passing an object by reference or by value.

## Arguments and Results in RMI

- ◆ Semantics of passing arguments for RMI in object-oriented languages needs to be defined.
- ◆ Argument and Result passing in Java RMI :  
When type of parameter is defined as remote interface, argument or result is passed as ROID (*by reference*).  
Other non-remote objects may be passed *by value* if they are *serializable*.
- ◆ Which objects can be accessed by RMI ?  
Any object can be accessed by RMI  
Distinguish between remote objects and local objects. (e.g. Java)  
Use interface definition language (IDL)
- ◆ Problem: migration/replication

## Dynamic Binding

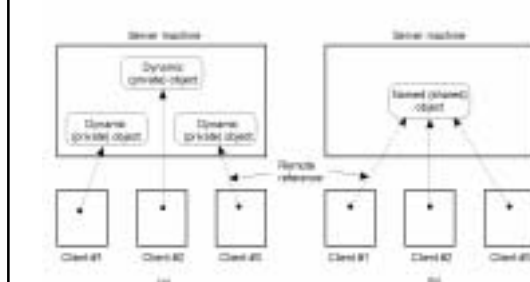
- ◆ Dynamic method binding should also apply to RMI.
- ◆ Smalltalk: Allow any message to be sent to any object, and raise exception if method is not supported.  
Distributed Smalltalk: general-purpose proxies.
- ◆ Java RMI :  
dynamic binding as a natural extension of local case  
Example:  

```
Shape aShape = (Shape) stack.pop();
float f = aShape.perimeter();
```

## Garbage Collection

- ◆ Some languages (Java, Smalltalk) support garbage collection.
- ◆ Explicit memory management difficult/impossible in distributed environment.
- ◆ Distributed garbage collection typically realized in ROID modules.  
Each ROID module:  
keeps track how many sites hold remote ROIDs for each local object  
(maintains **holders** table)  
informs other ROID modules about generation/deletion of ROIDs for their local objects (through the use of **addRef()** and **removeRef()**)
- ◆ Local garbage collector collects objects with no local or remote references.
- ◆ Reference counting (**addROID()** / **removeROID()**) over unreliable networks

## The DCE Distributed-Object Model



- Distributed dynamic objects in DCE.
- Distributed named objects

## Readings for Today's Lecture

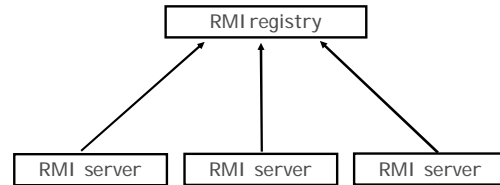
- References
  - Chapter 2 of "Distributed Systems: Principles and Paradigms"
  - Chapter 11 of "Java Network Programming and Distributed Systems"

## Java RMI

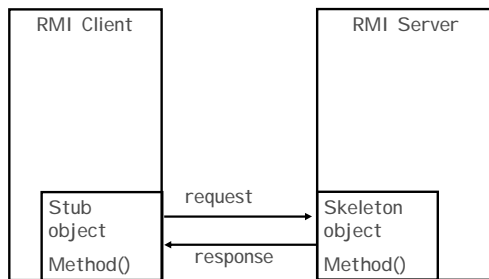
- RMI: A Java technology that allows one JVM to communicate with another JVM and have it execute an object method.
- RPC and RMI
  - RPC supports multiple languages, whereas RMI only support Java
  - RMI deals with objects, but RPC does not support the notion of objects
  - RPC offers procedures (not associated with a particular object)

## How RMI works

- ◆ The format used by RMI for representing a remote object reference: `rmi://hostname:port/serviceName`



## How RMI works



## Define a RMI Service Interface

```

Public interface RMI LightBulb extends java.rmi.Remote
{
    Public void on() throws java.rmi.remoteexecution;
    Public void off() throws java.rmi.remoteexecution;
    Public boolean ison() throws java.rmi.remoteexecution;
}
    
```

## Implement a RMI Service Interface

```

Public class RMI LightBulbImpl
    extends java.rmi.server.UnicastRemoteObject
    implements RMI LightBulb
{
    Public RMI LightBulbImpl () throws java.rmi.remoteexecution
    {setBulb(false);}

    Private boolean lighton;

    Public void on() throws java.rmi.remoteexecution
    {
        setBulb(true);
    }

    Public void off() throws java.rmi.remoteexecution
    {
        setBulb(false);
    }

    Public boolean ison() throws java.rmi.remoteexecution
    {return getBulb();}

    Public void setBulb(boolean value)
    {lighton = value;}

    Public void getBulb()
    {return lighton;}
}
    
```

## Create Stub and Skeleton Classes

Rmic RMI LightBulbImpl

Two files would be produced:

- RMI LightBulbImpl\_Stub.class
- RMI LightBulbImpl\_Skeleton.class

## Create a RMI Server

```
import java.rmi.*;
import java.rmi.server.*;

public class LightBulbServer
{
    public static void main(String args[])
    {
        try{
            RMI LightBulbImpl bulbService=new RMI LightBulbImpl();
            RemoteRef location = bulbService.getRef();

            String registry = args[0];

            String registration = "rmi://" + registry + "/RMI LightBulb";

            Naming.rebind(registration, bulbService);
        }
    }
}
```

## Create a RMI Client

```
import java.rmi.*;

public class LightBulbClient
{
    public static void main(String args[])
    {
        try{
            String registry = args[0];

            String registration = "rmi://" + registry + "/RMI LightBulb";

            Remote remoteService = Naming.lookup(registration);

            bulbService.on();
            system.out.println(bulbService.isOn());

            bulbService.off();
            system.out.println(bulbService.isOn());
        }
    }
}
```

## Running the RMI system

- ◆ Copy all necessary files to a directory on the local file system of all clients and the server.
- ◆ Change to the directory where the files are located, and run rmiregistry.
- ◆ In a separate console window, run the server with a hostname where rmiregistry is running.  
*Java LightBulbServer hostname*
- ◆ In a separate console window (another machine), run the client with a hostname where rmiregistry is running.  
*Java LightBulbServer hostname*

Any Questions?

See you next time.