

CprE 450/550x
Distributed Systems and Middleware

Inter-process Communication

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Potential Topics of Term Papers

The topic should be related to distributed systems in general, such as

- Group Communications,
- Peer-to-Peer systems,
- Overlay Networks,
- Grid Computing,
- Object Middleware,
- and others.

Potential Topics of Term Papers

- ◆ Resource Discovery and Management
- ◆ Security and Policy Management
- ◆ Resource Scheduling and Load Balancing
- ◆ Synchronization (e.g., clock synchronization, Election Algorithm, mutual exclusion, etc.)
- ◆ Consistency and replication
- ◆ Reliability and Survivability
- ◆ Performance Evaluation
- ◆ Anonymity. Censor-resistant
- ◆ Workload characterization
- ◆ Multi-cast Fingerprinting
- ◆ Anonymous authentication in dynamic group communications
- ◆ Data Replication strategies for Grid Computing systems
- ◆ Reputation-based resource scheduling for Grid Computing systems
- ◆ Middleware-based application design and development (e.g., real-time CORBA, fault-tolerant CORBA, etc.)
- ◆ You are welcome to propose your own topic!!!

Your Term Papers

- ◆ Every student is required to finish a term paper.
- ◆ Deadlines:
 - February 26 (Thursday, 5:00pm), topic selection due (UG).
 - March 11 (Thursday, 5:00pm), 3-pages proposal (problem definition) due (UG).
 - March 25 (Thursday, 5:00pm), 5-pages solutions and drafted evaluation plan due (G).
 - April 6 (Tuesday, 5:00pm), 5-pages summary of the papers you read and important issues you think (U)
 - April. 15 (Thursday 5:00pm), experimental results and improved solution due (G).
 - April. 27 (Tuesday 5:00pm), 10-page literature survey (U) and 15-pages full term paper (G) due, including problem definition, solutions, experimental data, conclusion and future works.
- ◆ Paper format: Latex or WORD, IEEE transactions, please refer: <http://www.ieee.org/organizations/pubs/transactions/stylesheets.htm>

Readings for Today's Lecture

- References
 - **Chapter 2 of "Distributed Systems: Principles and Paradigms"**

Object-Oriented Distributed Technology

- ◆ Objects
- ◆ Objects in Distributed Systems
- ◆ Requirements of Multi-User Applications

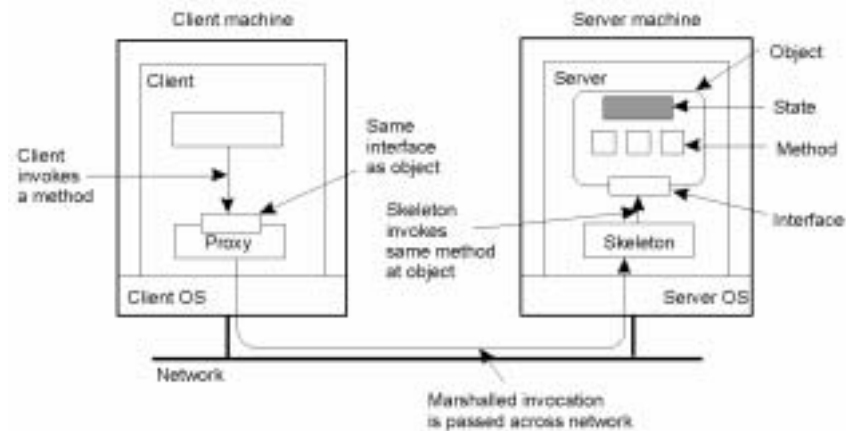
Object-Oriented Languages

- ◆ Object Identity
 - “object identifiers” (OIDs)
 - OIDs as first class values
- ◆ Actions
 - Initiated by sending message to object requesting method invocation
 - State in object may change
 - cascaded invocations of methods
- ◆ Dynamic Binding
 - The method executed is chosen according to the class of the recipient of the message.
- ◆ Garbage Collection
 - Dynamically allocated instances may be explicitly *deleted* or space is freed implicitly by garbage collector.

Objects in Distributed Systems

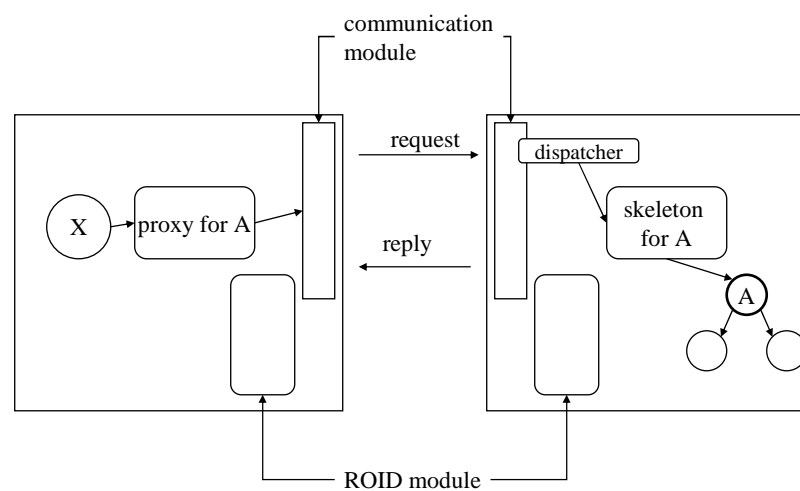
- ◆ Object Identity in a Distributed System
 - Remote object identifiers (ROIDs)
 - Ex. Java: ROID = endpoint (Java vm) + identifier (ObjID)
 - ROIDs as first-class values
 - Service for comparing remote object identifiers
 - e.g. Java: *RemoteObject::equals()*
- ◆ Actions in a Distributed Object System
 - Remote Method Invocation
- ◆ The Role of Proxies for Transparent RMI
 - Local proxy for each remote object that can be invoked by local object.
 - Local proxy behaves like local object, but, instead of executing message, forwards it to the remote object. (client stubs)
 - Remote object has skeleton object with server stub procedures

Distributed Objects



- ◆ Common organization of a remote object with client-side proxy.

Proxies and Skeletons



Proxies and Skeletons (cont)

◆ Proxies:

Need proxies to invoke remote objects.

Proxies are created when needed whenever ROID arrives in Reply message.

ROID module manages proxies and ROIDs.

◆ Dispatchers and Skeletons:

Not necessary for systems with reflection capabilities.

e.g. class *Method* in Java 1.2 reflection package:

method *invoke* can be called on instance of *Method*.

Dispatcher now generic and skeleton unnecessary.

Binding a Client to an Object

```
Distr_object* obj_ref;           //Declare a systemwide object reference
obj_ref = ...;                  // Initialize the reference to a distributed object
obj_ref-> do_something();        // Implicitly bind and invoke a method
```

(a)

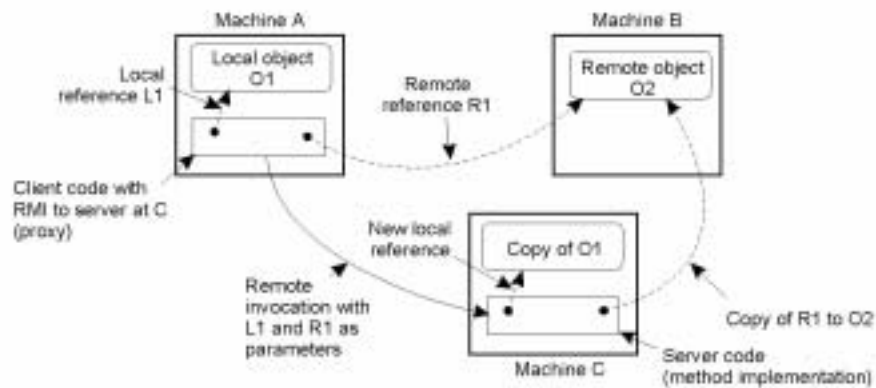
```
Distr_object objRef;           //Declare a systemwide object reference
Local_object* obj_ptr;         //Declare a pointer to local objects
obj_ref = ...;                 //Initialize the reference to a distributed object
obj_ptr = bind(obj_ref);       //Explicitly bind and obtain a pointer to the local proxy
obj_ptr -> do_something();      //Invoke a method on the local proxy
```

(b)

(a) Example with implicit binding using only global references

(b) Example with explicit binding using global and local references

Parameter Passing



- ◆ The situation when passing an object by reference or by value.

Arguments and Results in RMI

- ◆ Semantics of passing arguments for RMI in object-oriented languages needs to be defined.
- ◆ Argument and Result passing in Java RMI :
 - When type of parameter is defined as remote interface, argument or result is passed as ROID (*by reference*).
 - Other non-remote objects may be passed *by value* if they are *serializable*.
- ◆ Which objects can be accessed by RMI ?
 - Any object can be accessed by RMI
 - Distinguish between remote objects and local objects. (e.g. Java)
 - Use interface definition language (IDL)
- ◆ Problem: migration/replication

Dynamic Binding

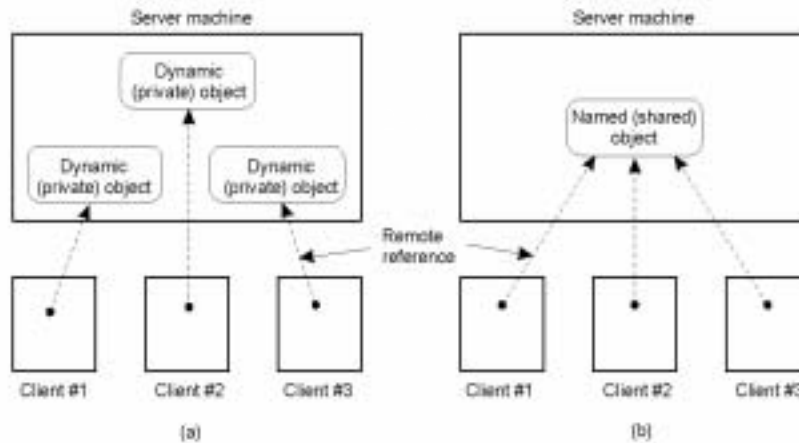
- ◆ Dynamic method binding should also apply to RMI .
- ◆ Smalltalk: Allow any message to be sent to any object, and raise exception if method is not supported.
Distributed Smalltalk: general-purpose proxies.
- ◆ Java RMI :
dynamic binding as a natural extension of local case
Example:

```
Shape aShape = (Shape) stack.pop();
float f = aShape.perimeter();
```

Garbage Collection

- ◆ Some languages (Java, Smalltalk) support garbage collection.
- ◆ Explicit memory management difficult/impossible in distributed environment.
- ◆ Distributed garbage collection typically realized in ROID modules.
Each ROID module:
 - keeps track how many sites hold remote ROIDs for each local object
(maintains **holders** table)
 - informs other ROID modules about generation/deletion of ROIDs for their local objects (through the use of **addRef()** and **removeRef()**)
- ◆ Local garbage collector collects objects with no local or remote references.
- ◆ Reference counting (**addROID()**/**removeROID()**) over unreliable networks

The DCE Distributed-Object Model



- a) Distributed dynamic objects in DCE.
 b) Distributed named objects

Readings for Today's Lecture

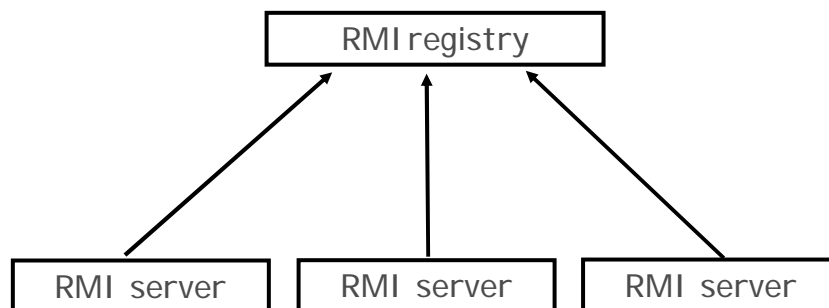
- References
 - Chapter 2 of "Distributed Systems: Principles and Paradigms"
 - Chapter 11 of "Java Network Programming and Distributed Systems"

Java RMI

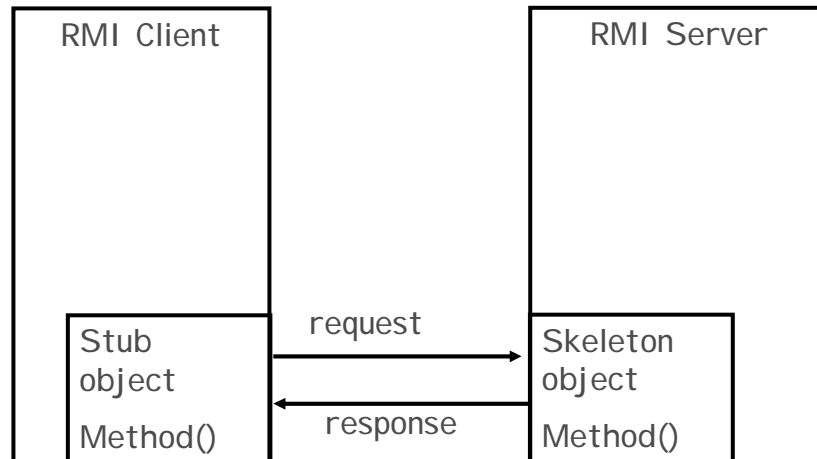
- RMI: A Java technology that allows one JVM to communicate with another JVM and have it execute an object method.
- RPC and RMI
 - RPC supports multiple languages, whereas RMI only support Java
 - RMI deals with objects, but RPC does not support the notion of objects
 - RPC offers procedures (not associated with a particular object)

How RMI works

- ◆ The format used by RMI for representing a remote object reference: `rmi://hostname:port/servicename`



How RMI works



Define a RMI Service Interface

```

Public interface RMI LightBulb extends java.rmi.Remote
{
    Public void on() throws java.rmi.remoteexecution;
    Public void off() throws java.rmi.remoteexecution;
    Public boolean ison() throws java.rmi.remoteexecution;
}
  
```

Implement a RMI Service Interface

```

Public class RMILightBulbImpl
    extends java.rmi.server.UnicastRemoteObject
    implements RMILightBulb
{
    Public RMILightBulbImpl() throws java.rmi.RemoteException
    {setBulb(false);}

    Private boolean lighton;

    Public void on() throws java.rmi.RemoteException
    {    setBulb(true); }

    Public void off() throws java.rmi.RemoteException
    {    setBulb(false); }

    Public boolean ison() throws java.rmi.RemoteException
    {return getBulb();}

    Public void setBulb(boolean value)
    {lighton = value;}

    Public void getBulb()
    {return lighton;}

}

```

Create Stub and Skeleton Classes

`Rmic RMILightBulbImpl`

Two files would be produced:

- `RMILightBulbImpl_Stub.class`
- `RMILightBulbImpl_Skeleton.class`

Create a RMI Server

```
import java.rmi.*;
import java.rmi.server.*;

Public class LightBulbServer
{
    Public static void main(String args[])
    {
        Try{
            RMI LightBulbI mpl bulbService=new RMI LightBulbI mpl();
            RemoteRef location = bulbService.getRef();

            String registry = args[0];

            String registration = "rmi://" + registry + "/RMI LightBulb";

            Naming.rebind(registration, bulbService);

        }
    }
}
```

Create a RMI Client

```
import java.rmi.*;

Public class LightBulbClient
{
    Public static void main(String args[])
    {
        Try{
            String registry = args[0];

            String registration = "rmi://" + registry + "/RMI LightBulb";

            Remote remoteService = Naming.lookup(registration);

            bulbService.on();
            system.out.println(bulbService.isOn());

            bulbService.off();
            system.out.println(bulbService.isOn());

        }
    }
}
```

Running the RMI system

- ◆ Copy all necessary files to a directory on the local file system of all clients and the server.
- ◆ Change to the directory where the files are located, and run `rmiregistry`.
- ◆ In a separate console window, run the server with a hostname where `rmiregistry` is running.

Java LightBulbServer hostname

- ◆ In a separate console window (another machine), run the client with a hostname where `rmiregistry` is running.

Java LightBulbServer hostname

Any Questions?

See you next time.